

A PageRank Algorithm based on Asynchronous Gauss-Seidel Iterations

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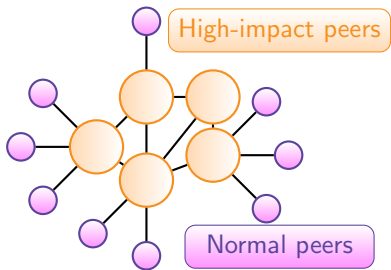
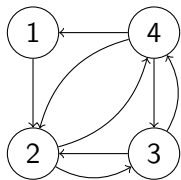
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- 2 Problem Statement
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- 4 Results
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Motivation

- Traditional PageRank - Compute relative importance of each webpage in the network.
- Ranking of Researchers - PageRank has also been proposed to be an alternative to citation count.
- General Ranking - Since PageRank is equivalent to the relative time spent in a node for a random walk on the graph many more applications are possible.



Intuition behind PageRank

- The ranks should translate to the relative time one would spent on that node for a random walk.
- It should be a steady-state vector.
- Dangling nodes must not appear observant and the random walk must be able to jump to different partitions.
- Problem: $x^* \in [0, 1]^n$, $\sum_{i=1}^n x_i^* = 1$

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- Problem:

$$x^* = Mx^*, \quad x^* \in [0, 1]^n, \quad \sum_{i=1}^n x_i^* = 1$$

$$M = (1 - m)A + \frac{m}{n}S, \quad m = 0.15, \quad S := \mathbf{1}_n \mathbf{1}_n^T$$

Some related work

- Defining distributed randomized versions of the power method:
[C. Ravazzi et al. \(2015\)](#). “Ergodic Randomized Algorithms and Dynamics Over Networks”. In: *IEEE Transactions on Control of Network Systems*
- Use of sequential Gauss-Seidel with no projection:
[A. Arasu et al. \(2002\)](#). “PageRank computation and the structure of the Web: experiments and algorithms”. In: *The Eleventh International WWW Conference*. ACM Press
- Using stochastic approximation of the gradient:
[J. Lei and H. F. Chen \(2015\)](#). “Distributed Randomized PageRank Algorithm Based on Stochastic Approximation”. In: *IEEE Transactions on Automatic Control*

Equivalent formulations for PageRank

- Using the power method it becomes:

$$x(k+1) = Mx(k) = (1-m)Ax(k) + \frac{m}{n}\mathbf{1}_n$$

- Intuitively, the Power Method searches for the eigenvector;
- It can be solved as an optimization problem;
- The rank vector is also a eigenvector:

$$(I_n - (1-m)A)x = \frac{m}{n}\mathbf{1}_n.$$

Problem Statement

- Find an algorithm that can exploit the fact that subsets of pages are stored in the same processor;
- The iteration should involve A and not M to keep sparsity;
- Solve the equation:

$$(I_n - (1 - m)A)x = \frac{m}{n} \mathbf{1}_n \quad (1)$$

with constraints:

$$x^* \in [0, 1]^n, \quad \sum_{i=1}^n x_i^* = 1 \quad (2)$$

PageRank Algorithm Problem

How to find a faster algorithm satisfying the above restrictions?

Power Method as the Jacobi Method

- For a linear equation $Ax = b$ partition $A = D + R$ where $D = \text{diag}(A)$ and $R = A - \text{diag}(A)$;
- The Jacobi Method is given by the iteration:

$$x(k+1) = D^{-1}(b - Rx(k)) \quad (3)$$

- If we take $D = I_n$, $R = -(1 - m)A$ and $b = \frac{m}{n} \mathbf{1}_n$ the Jacobi Method is equivalent to the Power Method;
- Jacobi Method is typically used for solving diagonally dominant equations.

PageRank using Gauss-Seidel

- Partitioning $A = L + D + U$, for lower and upper triangular matrices L and U and diagonal D ;
- The Gauss-Seidel Method becomes:

$$x(k + 1) = (L + D)^{-1}(b - Ux(k)) \quad (4)$$

- The method can be written to take advantage of updated values

$$x_i(k + 1) = \frac{1}{A_{ii}} \left(b_i - \sum_{j=1}^{i-1} A_{ij}x_j(k + 1) - \sum_{j=i+1}^n A_{ij}x_j(k) \right). \quad (5)$$

- The Gauss-Seidel is also an iterative algorithm for linear equations but corresponding to a different partitioning of the matrix.

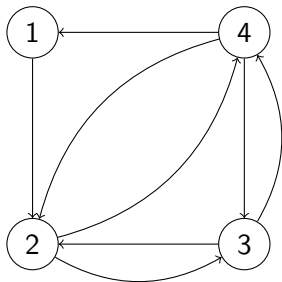
Results

- Theoretical result: $\rho(T_{gs}) < \rho(T_p)$
- Main challenge: computed x^* using the Gauss-Seidel might not sum to one.
- It is simulated:
 - Projection on the n -simplex;
 - Normalization;
 - Randomized and Asynchronous versions.

Simulation Results (1/2)

Setup: Simple case of four pages.

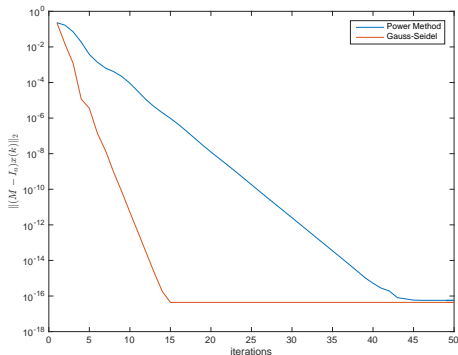
- Gauss-Seidel is faster than the Power Method.
- Performance is affected with no projection.
- Normalization and projection are similar.
- Asynchronous version has an intermediary performance.



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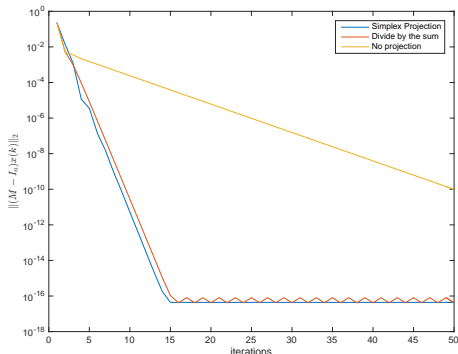
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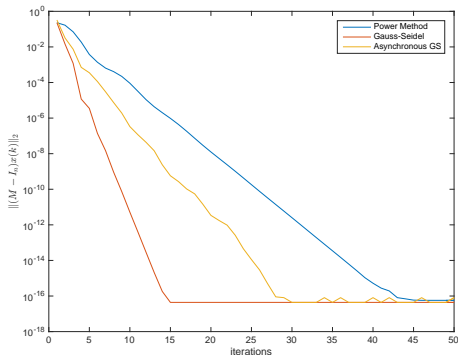
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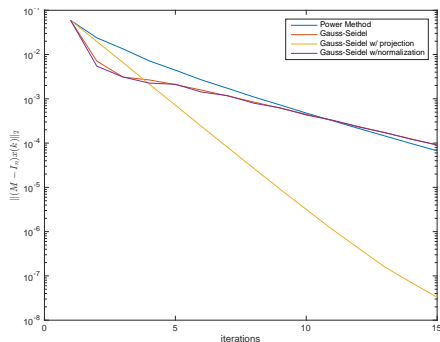
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Simulation Results (2/2)

Setup: 1000 Monte Carlo experiment for a 500-page network generated by Barabási-Albert algorithm to simulate a power law network as the World-Wide Web.

- In a more realistic scenario the projection is key to get very small errors;
- Fast initial convergence with normalization;
- The problem of having Gauss-Seidel going to zero can be avoided by normalization.



Concluding Remarks

Contributions:

- We have shown that Gauss-Seidel is faster than the Power Method
- Gauss-Seidel does not keep the sum of the state
- Two approaches:
 - projection - implies two communication rounds;
 - normalization - no added complexity.
- Shown how it can be made Asynchronous and Randomized with performance between the Power Method and the Synchronous Gauss-Seidel.

The end

- Thank you for your time.

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